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## USING RATIONAL NUMBERS AND PARALLEL COMPUTING TO EFFICIENTLY AVOID ROUND-OFF ERRORS ON MAP SIMPLIFICATION

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Maurício G. Gruppi<sup>1</sup>

Salles V. G. de Magalhães<sup>1,2</sup>

Marcus V. A. Andrade<sup>1</sup>

W. Randolph Franklin<sup>2</sup>

Wenli Li<sup>2</sup>

<sup>1</sup>Departamento de Informática - Universidade Federal de Viçosa

<sup>2</sup>Rensselaer Polytechnic Institute - USA



## 1. Introduction

- What is Map Simplification?
- Topological Consistency

## 2. Related Works

- Simplification Algorithms
- Round-off Errors in Floating-point Arithmetic

## 3. Round-off Errors on Map Simplification

## 4. The EPLSimp Method

## 5. Experimental Evaluation

## 6. Conclusions



## INTRODUCTION

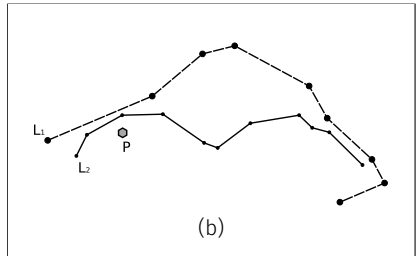
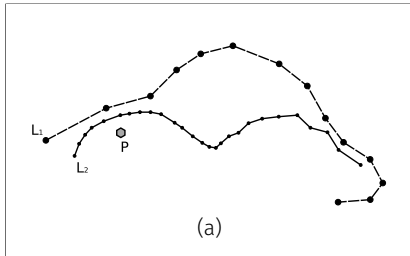
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## What is Map Simplification?

- It's the process of reducing the amount of detail of a map.
- Such as reducing the number of vertices of a **polygonal chain** when altering scale.
- However, there are key features that must be preserved.



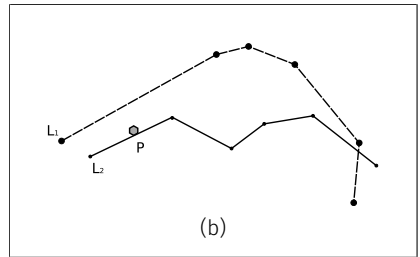
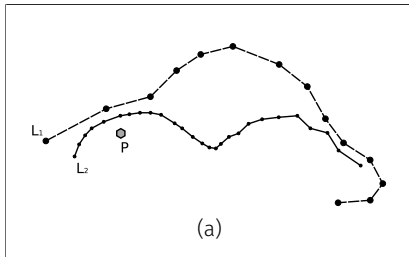
# MAP SIMPLIFICATION



(a) shows an example of input set. A **topologically consistent** simplification of (a) is shown in (b).



# MAP SIMPLIFICATION



(b) shows a **topologically inconsistent** simplification of (a).



## RELATED WORKS

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## Ramer-Douglas-Peucker's Algorithm (RDP)

[Douglas and Peucker, 1973][Ramer, 1972]

- Simplification by *selection*.
- May produce inconsistency.
  - [Saalfeld, 1999]
  - [Li et al., 2013]





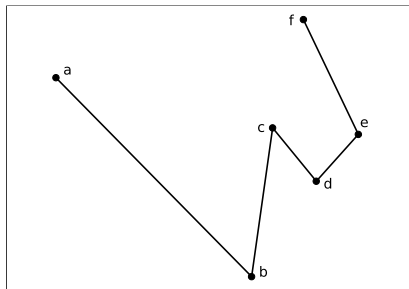
## Visvalingam-Whyatt's Algorithm (VW) [Visvalingam and Whyatt, 1993]

- Simplification by *elimination*.
- Ranks points by *effective area*.
- Removes points whose *effective area* is smaller than a given threshold.



# SIMPLIFICATION ALGORITHMS

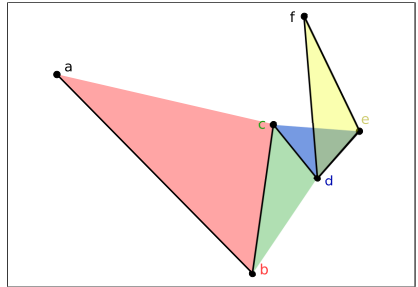
## VW's Algorithm



# SIMPLIFICATION ALGORITHMS

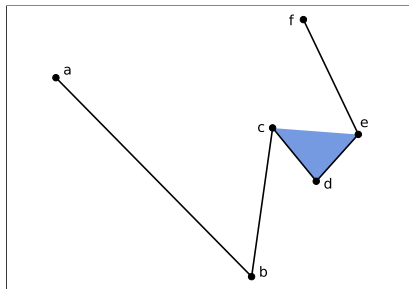
## VW's Algorithm

- Calculates every point's *effective area* in  $L$ .
- **Definition:** The *effective area* of a polyline vertex  $v_i$  is the area of the triangle formed by  $v_{i-1}$ ,  $v_i$  and  $v_{i+1}$ .



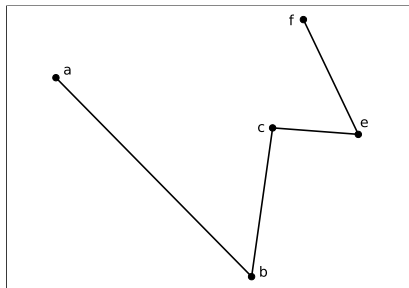
## VW's Algorithm

- Calculates every point's *effective area* in  $L$ .
- Find the point  $p$  with smallest effective area.



## VW's Algorithm

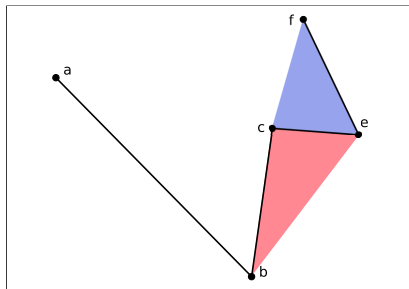
- Calculates every point's *effective area* in  $L$ .
- Find the point  $p$  with smallest effective area.
- Remove  $p$  from  $L$ .



# SIMPLIFICATION ALGORITHMS

## VW's Algorithm

- Calculates every point's *effective area* in  $L$ .
- Find the point  $p$  with smallest effective area.
- Remove  $p$  from  $L$ .
- Calculate the new effective area for  $p$ 's neighbors.



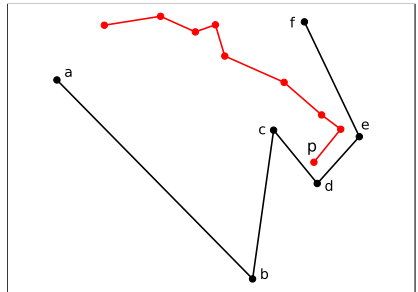
VW's algorithm can produce topologically inconsistent results.



# SIMPLIFICATION ALGORITHMS

## VW's Algorithm

- Input: 2 polygonal chains (**black** and **red**).

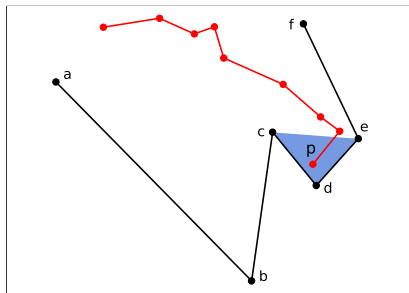




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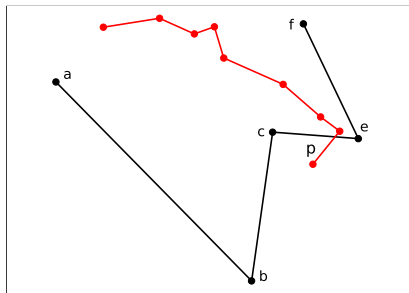
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- Input: 2 polygonal chains (**black** and **red**).
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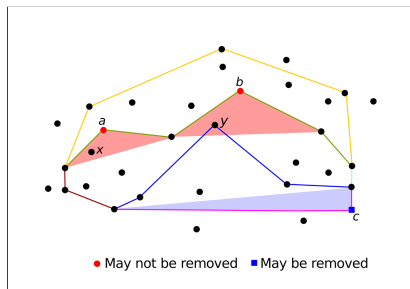
## VW's Algorithm

- Input: 2 polygonal chains (**black** and **red**).
- Remove  $d$  from  $L$ .
- An intersection has been created between both lines.



## TopoVW [Gruppi et al., 2015]

- Ranks points by effective area.
- Checks for points inside each  $p$ 's triangle.
- Removes  $p$  if there is none.
- Stops when a certain number of points have been removed.



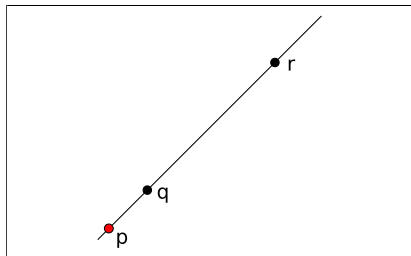
# ROUND-OFF ERRORS IN FLOATING POINT ARITHMETIC

- Algorithms previously mentioned were designed for floating-point arithmetic.
- Arbitrary precision numbers represented as fixed precision numbers.
- May incur round-off errors.
- Therefore producing **wrong** results.



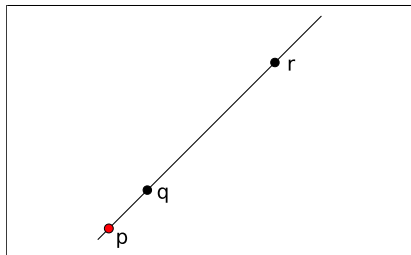
# ROUND-OFF ERRORS IN FLOATING POINT ARITHMETIC

- Round-off errors affect planar orientation predicate [Kettner et al., 2008].
- The problem of finding whether three points  $p$ ,  $q$ ,  $r$ :
  - are collinear.
  - make a left-turn.
  - make a right-turn.



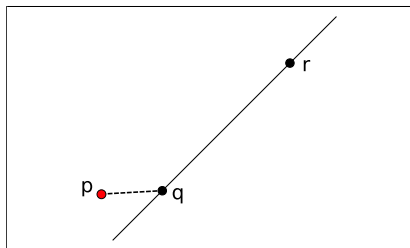
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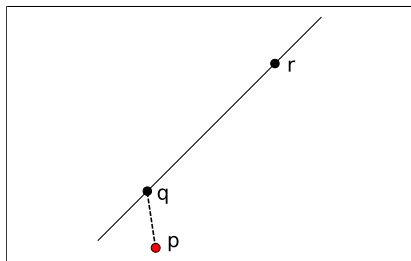
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# ROUND-OFF ERRORS IN FLOATING POINT ARITHMETIC

$$orientation = \text{sign} \left( \begin{vmatrix} 1 & p_x & p_y \\ 1 & q_x & q_y \\ 1 & r_x & r_y \end{vmatrix} \right)$$



# ROUND-OFF ERRORS IN FLOATING POINT ARITHMETIC

$$\text{orientation} = \text{sign} \left( \begin{vmatrix} 1 & p_x & p_y \\ 1 & q_x & q_y \\ 1 & r_x & r_y \end{vmatrix} \right)$$

Sign:

- +: left turn.
- -: right turn.
- 0: collinear.



# ROUND-OFF ERRORS IN FLOATING POINT ARITHMETIC

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**Sign:**

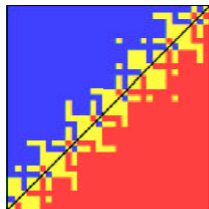
- +: left turn.
- -: right turn.
- 0: collinear.

**Possible problems:**

- *rounding to zero.*
- *perturbed zero.*
- *sign-inversion.*



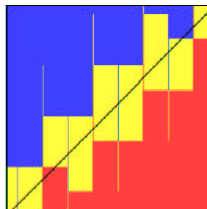
# ROUND-OFF ERRORS IN FLOATING POINT ARITHMETIC



$$p: \begin{pmatrix} 0.5 \\ 0.5 \end{pmatrix}$$

$$q: \begin{pmatrix} 12 \\ 12 \end{pmatrix}$$

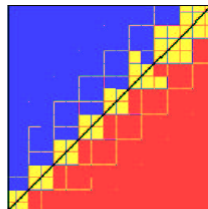
$$r: \begin{pmatrix} 24 \\ 24 \end{pmatrix}$$



$$\begin{pmatrix} 0.500000000000002531 \\ 0.50000000000000171 \end{pmatrix}$$

$$\begin{pmatrix} 17.300000000000001 \\ 17.300000000000001 \end{pmatrix}$$

$$\begin{pmatrix} 24.000000000000005 \\ 24.0000000000000517765 \end{pmatrix}$$



$$\begin{pmatrix} 0.5 \\ 0.5 \end{pmatrix}$$

$$\begin{pmatrix} 8.8000000000000007 \\ 8.8000000000000007 \end{pmatrix}$$

$$\begin{pmatrix} 12.1 \\ 12.1 \end{pmatrix}$$

Result of the planar orientation problem using floating-point arithmetic.  
Source: [Kettner et al., 2008].



## ROUND-OFF ERRORS ON MAP SIMPLIFICATION

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# ROUND-OFF ERRORS IN MAP SIMPLIFICATION

We have tested floating-point round-off errors on map simplification:

- We needed to determine whether  $p$  is inside triangle  $T$  formed by  $(r, s, t)$ .
- This was done by using *barycentric coordinates* of  $p$  in  $T$ .



# ROUND-OFF ERRORS ON MAP SIMPLIFICATION

Let  $a$ ,  $b$  and  $c$  be scalars such that:

- $p_x = ar_x + bs_x + cr_x$
- $p_y = ar_y + bs_y + cr_y$
- $a + b + c = 1$

$p$  lies inside  $T$  if and only if  $0 \leq a, b, c \leq 1$



# ROUND-OFF ERRORS ON MAP SIMPLIFICATION

- A function *is\_inside*( $r, s, t, p$ ) was implemented in C++ using floating-point numbers.





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# ROUND-OFF ERRORS ON MAP SIMPLIFICATION

- A function *is\_inside*( $r, s, t, p$ ) was implemented in C++ using floating-point numbers.
- *false inside*: outer point said inside.
- May prevent simplification.



# ROUND-OFF ERRORS ON MAP SIMPLIFICATION

- A function *is\_inside*( $r, s, t, p$ ) was implemented in C++ using floating-point numbers.
- *false inside*: outer point said inside.
- *false outside*: inner point said outside.

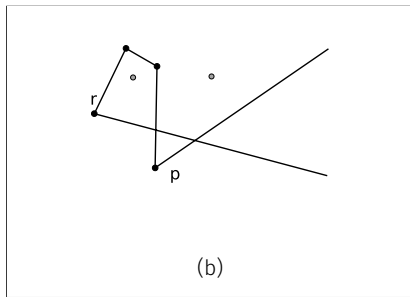
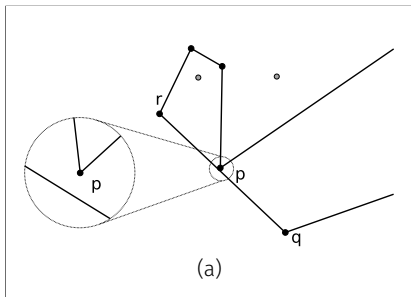


# ROUND-OFF ERRORS ON MAP SIMPLIFICATION

- A function *is\_inside*( $r, s, t, p$ ) was implemented in C++ using floating-point numbers.
- *false inside*: outer point said inside.
- *false outside*: inner point said outside.
- May create improper intersections and self-intersections.



# ROUND-OFF ERRORS ON MAP SIMPLIFICATION



$p$  was a **false outside**. Thus the removal of  $q$  was possible, creating self-intersections.



# SOLUTION TO ROUND-OFF ERRORS

- $\epsilon$ -tweaking
- Snap-rounding
- Exact Arithmetic



$\epsilon$ -*tweaking* uses a tolerance value when comparing two numbers:

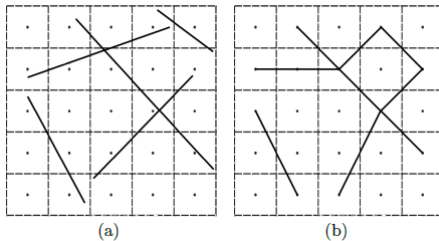
$$x = y, \text{ if } |x - y| \leq \epsilon.$$

- Automatically activates *rounding to zero*.
- Finding  $\epsilon$  is difficult. Especially for big datasets.



# SNAP-ROUNDING

*Snap-rounding* splits the map into pixels (cells). Rounds every endpoint to the center of its bounding pixel.



**Figure:** (a) before snap-rounding. (b) after snap-rounding. Intersections were introduced.





# EXACT ARITHMETIC WITH RATIONAL NUMBERS

*Exact Arithmetic with Rational Numbers:*

- Non-integer variables are represented as arbitrary precision rational numbers.
- Slower than floating-point arithmetic but round-off errors free.
- Overhead can be reduced using parallel computation.

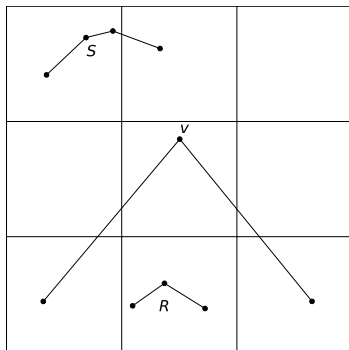


- **EPLSimp** uses exact arithmetic for simplifying polylines.
- A uniform-grid structure is used for determining which points needed to be tested for each triangle.
- Parallel computing used for performance.
- Lines are then subdivided into sets that can be simplified in parallel.



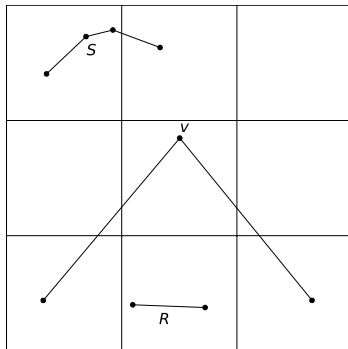
# OUR METHOD

- Construct a uniform-grid in parallel.



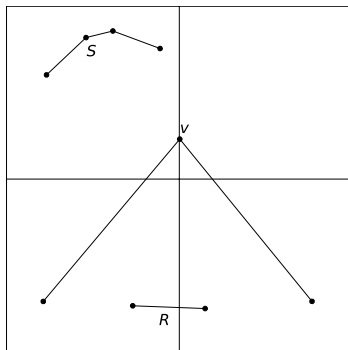
# OUR METHOD

- Construct a uniform-grid in parallel.
- Simplify line  $R$ .

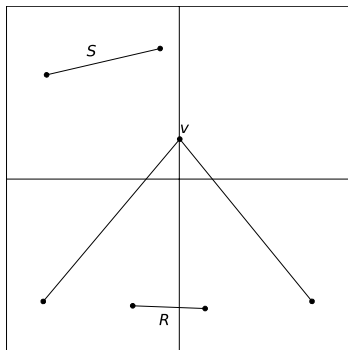


# OUR METHOD

- Construct a uniform-grid in parallel.
- Simplify line  $R$ .
- Decrease grid resolution.
- More lines inside single cells.
- Allows parallel simplification.

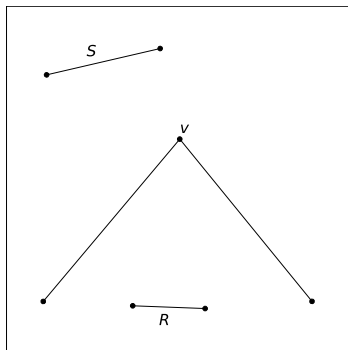


- Simplify line S.

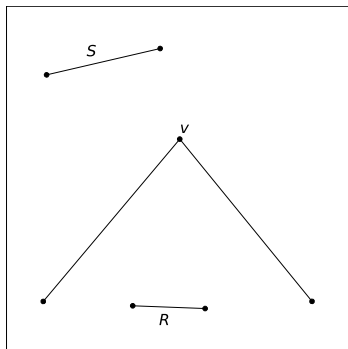


# OUR METHOD

- Simplify line S.
- Decrease grid resolution.



- Simplify line  $S$ .
- Decrease grid resolution.
- Simplify the remaining lines (if any).





## EXPERIMENTAL RESULTS

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# EXPERIMENTAL RESULTS - CONSISTENCY

- *EPLSimp* was implemented in C++ using the GMPXX library [Granlund and the GMP development team, 2014].
- Artificial datasets were created to evaluate the occurrence of round-off errors.
- *EPLSimp* did not produce any topological inconsistencies.



# EXPERIMENTAL RESULTS - PERFORMANCE

**Table:** Times (in ms) for the main steps of the map simplification algorithms. Rows *Max* represent the time for removing the maximum amount of points from the map while rows *Half* represent the time to remove half of the points.

	Dataset Method	1		2		3	
		TopoVW	<i>EPLSimp</i>	TopoVW	<i>EPLSimp</i>	TopoVW	<i>EPLSimp</i>
Max.	Initialize	4	22	28	190	1828	5353
	Simplify	39	60	626	445	46069	57095
	Total	43	82	654	635	47897	62448
Half	Initialize	4	22	28	186	1847	5447
	Simplify	25	41	357	331	23021	48090
	Total	29	63	384	517	24868	53537



# EXPERIMENTAL RESULTS - PERFORMANCE

**Table:** Times (in ms) for initializing and simplifying maps from the 3 datasets considering different number of threads. The simplification was configured to remove the maximum amount of points from the maps.

		Initialization			Simplification		
Dataset		1	2	3	1	2	3
Threads	1	71	655	26833	176	1574	250237
	2	91	568	15483	152	1150	131310
	4	54	422	9853	99	689	82641
	8	34	240	6552	61	483	62089
	16	22	190	5353	60	445	57095



- We were able to avoid round-off errors using exact arithmetic with rational numbers.
- Parallel computing helped alleviating the overhead, approaching floating-point's processing time.
- Future works include:
  - Adapting *EPLSimp* for simplifying vector drawings and 3D objects.
  - Use exact arithmetic for other GIS algorithms.



# THANK YOU



mauricio.gruppi@ufv.br  
salles@ufv.br  
marcus@ufv.br  
mail@wrfranklin.org