

PROB 2(3|22 - 1

POISSON DISTRIBUTION

of RADIOACTIVE DECAY IN A BLOCK OF RADIUM,

THERE ARE MANY ATOMS, EACH VERY UNLUCKY TO DECAY IN NEXT SECOND.

BUT SOME WILL.

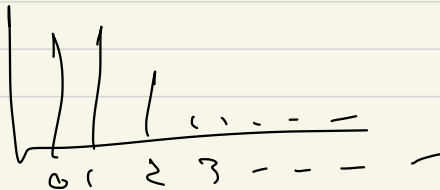
α = EXPECTED NUMBER TO DECAY IN NEXT SECOND

$$P(K \text{ WILL DECAY}) = \frac{\alpha^k}{k!} e^{-\alpha}$$

$$\alpha = 1. \quad P(k) = \frac{1^k}{k!} e^{-1} \approx \frac{.3}{k!}$$

$$P(0) = .3 \quad P(2) = .05$$

$$P(1) = .3 \quad P(3) = .05 \quad \dots$$



$$\alpha = 2$$

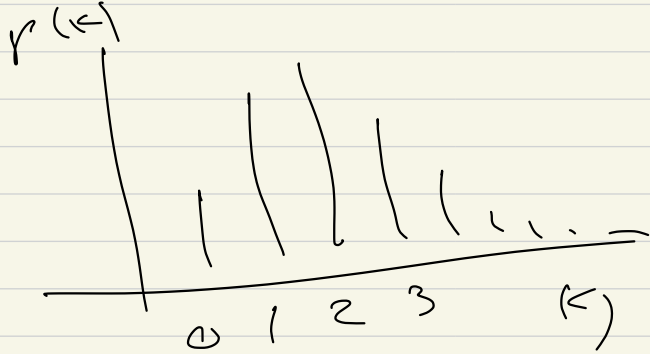
$$p(k) = \frac{2^k}{k!} e^{-2} \sim \frac{2^k}{k!}$$

$$p(0) = 1$$

$$p(1) = 2$$

$$p(2) = 2$$

$$p(3) = 1.25$$



$$p(k) = \frac{\alpha^k}{k!} e^{-\alpha}$$

CHECK $\sum_{k=0}^{\infty} p(k) = \sum_{k=0}^{\infty} \frac{\alpha^k}{k!} e^{-\alpha}$

$$\sum_{k=0}^{\infty} \frac{\alpha^k}{k!} = e^{\alpha}$$

$$E[k] = \sum_{k=p}^{\infty} k p(k) = \sum_{k=p}^{\infty} k \frac{\alpha^k}{k!} e^{-\alpha}$$

$$= \alpha e^{-\alpha} \sum_{k=p}^{\infty} \frac{1}{(k-1)!} \alpha^{k-1}$$

I CAN CHANGE LOWER LIMIT FROM 0 TO 1 BECAUSE $k p(k) = 0$ WHEN $k=0$

$$\sum_{k=0}^{\infty} k p(k) = \sum_{k=1}^{\infty} k p(k)$$

$$E[k] = \alpha e^{-\alpha} \sum_{k=1}^{\infty} \frac{1}{(k-1)!} \alpha^{k-1}$$

$$= \alpha e^{-\alpha} \left(\sum_{k=0}^{\infty} \frac{\alpha^k}{k!} \right) \leftarrow e^{\alpha}$$

REPLACE $(k-1)$ BY k EVERYWHERE.

$$E[k] = \alpha e^{-\alpha} e^{\alpha} = \alpha$$

FOR VAR[K] FOR POISSON
NEED $E[K^2]$

$$= \sum_{k=1}^{\infty} k^2 \frac{\alpha^k}{k!} e^{-\alpha}$$

$$= e^{-\alpha} \left[\sum (k + k(k-1)) \frac{\alpha^k}{k!} \right]$$

$$= e^{-\alpha} \left[\sum \frac{k\alpha^k}{k!} + \sum \frac{k(k-1)\alpha^k}{k!} \right]$$

$$= e^{-\alpha} \left[\alpha \sum \frac{\alpha^{k-1}}{(k-1)!} + \alpha^2 \sum \frac{\alpha^{k-2}}{(k-2)!} \right]$$

$$= e^{-\alpha} [\alpha e^{\alpha} + \alpha^2 e^{\alpha}]$$

$$= \alpha + \alpha^2$$

$$E[K^2] = \alpha + \alpha^2$$

MUST CHECK BOUNDARIES

$$\text{VAR}[K] = E[K^2] - E[K]^2$$

$$\alpha + \alpha^2 - \alpha^2 = \alpha$$

FOR POISSON

$$E[K] = \text{VAR}(K) = \alpha$$

$$\sigma[K] = \sqrt{\alpha}$$



FOR BIGGER α , R.V. IS MORE TIGHTLY CLUSTERED, AS A PERCENTAGE.

$$\text{IF } \alpha = 1000 \quad \begin{array}{cc} 970 - 1030 \\ E - S & E + S \end{array}$$

ABOUT 2/3 OF TIME-

COMMUNICATION EXAMPLE.

$N = 48$ CUSTOMERS FOR A SHARED CHANNEL.

$\alpha = \frac{1}{3}$ OF THEM TALK AT ANY GIVEN TIME.

Q: HOW MANY CHANNELS SHOULD COMPANY PROVIDE?

(IF CUSTOMERS ARE INDEPENDENT LET R.V. $K = \#$ TALKING NOW)

$$P(K) = \binom{48}{k} \alpha^k (1-\alpha)^{48-k}$$

BINOMIAL

MEAN?

$$E[K] = \sum_{k=0}^{48} k \alpha^k (1-\alpha)^{48-k}$$

$$= 48\alpha = \underline{16}$$

ASIDE

BINOMIAL IS EXACT

(IF N IS VERY LARGE AND α

VERY SMALL AND k SMALL

THEN $P(k) = \binom{N}{k} \alpha^k (1-\alpha)^{N-k}$

IS HARD TO COMPUTE

(IF $E(k)$ IS SMALL NUMBER THEN

POISSON IS A GOOD APPROX.

(HERE α IS DIFFERENT)

ANOTHER APPROX IS

GAUSSIAN (NORMAL) DIST.

(IF $N = 6 \times 10^{23}$

$k = 5$

$\binom{6 \times 10^{23}}{5}$

IS A
MESS

APPROX FOR $N!$ = $\frac{N^N}{e^N} \sqrt{2\pi N}$

STIRLING'S
APPROX

8

TRY $N=2$ $\frac{2^2}{e^2} \sqrt{4\pi}$ $\frac{2\sqrt{\pi}}{e^2}$

EXACT IS 2.

RELATIVE ERROR $\frac{1}{(2N)}$ 1.919
NOT BAD

BACK TO COMMUNICATION EX.
 48 CUSTOMERS TALK $\frac{1}{3}$ OF TIME
 M CHANNELS.

HOW MANY UNHAPPY CUSTOMERS?
 IF $M=48$ NONE.

R.V $K = \#$ TALKING CUSTOMERS.

IF $K \leq M$ ALL HAPPY.

IF $K > M$ $K-M$ UNHAPPY.

LET $U = \#$ UNHAPPY.

$$U = (K-M)^+ \leftarrow \text{CLIPPED AT 0}$$

$$X^+ = \begin{cases} X & \text{if } X \geq 0 \\ 0 & \text{if } X < 0 \end{cases}$$

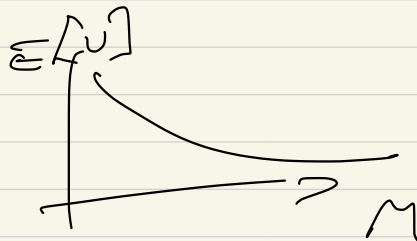
WHAT $E[U]$

$$\sum_{k=0}^{K-M} \binom{K-M}{k} \alpha^k (1-\alpha)^{K-M-k}$$

$$E[U] = \sum_{k=M+1}^K \binom{K-M}{k} \alpha^k (1-\alpha)^{K-k}$$

NO SIMPLE FORM.

YOU CAN PLOT



CONDITIONAL PMF

PMF = PROBABILITY MASS FN.

$$P(A|B) = \frac{P(A \cap B)}{P(B)}$$

X IS R.V. UNIFORM IN $[0, N-1]$

$$P(k) = \begin{cases} \frac{1}{N} & 0 \leq k \leq N-1 \\ 0 & \text{OTHERWISE} \end{cases}$$

k IS AN INTEGER.

SAY $N=10$ $P(k) = \frac{1}{10}$

SAY WE KNOW THAT $X \geq 5$

WHAT IS $P(X | X \geq 5)$?

$$\frac{P(X \cap X \geq 5)}{P(X \geq 5)}$$

0 1 2 3 4 | 5 6 7 8 9

IF $X < 5$ 0 }
IF $X \geq 5$

$$\frac{\frac{1}{10}}{\frac{2}{10}} = \frac{1}{5}$$

UNIFORM $[0 \dots N-1]$

N CHOICES.

K IS A POSSIBLE RESULT

$$P(K) = \begin{cases} \frac{1}{N} & \text{IF } 0 \leq K < N \\ 0 & \text{OTHERWISE} \end{cases}$$

TOSS COIN 3 TIMES

$K = \#$ HEADS.

IF WE KNOW $K \geq 1$, WHAT ARE PROBS FOR $K = 1, 2, 3$?

$$P(K | K \geq 1) = \frac{P(K \cap K \geq 1)}{P(K \geq 1)}$$

K

0

$$\frac{3}{8} = \frac{3}{7}$$

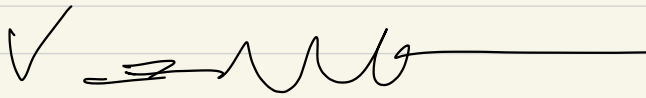
2 - $3(2/8) = 3/4$

3 - $1(2/8) = 1/4$

$P(K \geq 1)$

$\frac{7}{8}$

SQUARE LAW DEVICE



RESISTOR

1Ω

V IS R.V.

$\begin{matrix} -3 \\ -1 \\ 1 \end{matrix}$

$P = \frac{I^2}{4}$
EACH

CASE -

$$E[V] = \frac{1}{4} (-3 + -1 + 1 + \dots) = 0$$

$$\text{POWER} = V^2$$

$$E[V^2] = \sum \left[\frac{1}{4} \cdot 9 + \frac{1}{4} \cdot 1 + \frac{1}{4} \cdot 1 + \frac{1}{4} \cdot 9 \right]$$

$$= \frac{20}{4} = 5$$

TABLE 115-116 IS A USEFUL TABLE
OF DISCRETE R.V.